Database query languages

### on views

what are views usefull for?



# outline

- 1. what is a view?
- 2. processing queries involving views
- 3. what is a materialized view?
- 4. processing queries using materialized views
- 5. recursive views (see e.g., notes of CS145 from Stanford U.)

#### views

Movies[title,year,length,genre,studioName,producer] Producer[name,nationality]

CREATE VIEW MovieProd(movieTitle,prodNat) AS SELECT title,nationality FROM Movies, Producer WHERE producer = name

put differently MovieProd(t,n)  $\leftarrow$  Movies(t,y,g,s,p),Producer(p,n)

# querying views

- a view is not a table
- it is an intensional definition
- its extension is not stored
- but it may be queried as a table

SELECT starName,ProdNat FROM MovieProd, StarsIn WHERE title=movieTitle

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### processing queries using views

SELECT V.x FROM V WHERE  $\varphi$  (i.e.,  $\pi_{V.x}(\sigma_{\varphi}(V))$ )

suppose V is defined by:  $\pi_{y,x}(\sigma_{\theta}(R))$ 

processing Q means processing:  $\pi_x(\sigma_{\varphi}(\pi_{y,x}(\sigma_{\theta}(R))))$ 

and thus:  $\pi_x(\sigma_{\varphi \land \theta}(R)))$ 

materialized views

#### CREATE MATERIALIZED VIEW MovieProd(movieTitle,prodNat) AS SELECT title,nationality FROM Movies, Producer WHERE producer = name

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### materialized views

- a materialized view is not a table
- even though its extension is stored
- it remains an intensional definition
- and it may be queried as a table

if tuples are INSERTed (resp. DELETEd) in (resp. FROM) e.g., Movies, then MovieProd has to be refreshed

#### views maintenance

for simple views, maintenance can be incremental

- ► consider MovieProd defined as MovieProd(t,n) ← Movies(t,y,g,s,p),Producer(p,n)
- suppose (a,b,c,d,e) is inserted into Movies
- what to insert into MovieProd?
- we have the title:  $\pi_{title}(a,b,c,d,e)$
- SELECT nationality FROM Producer WHERE name='e'

#### views maintenance

for simple views, maintenance can be incremental

- ► consider MovieProd defined as MovieProd(t,n) ← Movies(t,y,g,s,p),Producer(p,n)
- suppose (a,b) is inserted into Producer
- what to insert into MovieProd?
- INSERT INTO MovieProd SELECT title, 'b' FROM Movies, Producer WHERE name='a'

#### views maintenance

for simple views, maintenance can be incremental

- ► consider MovieProd defined as MovieProd(t,n) ← Movies(t,y,g,s,p),Producer(p,n)
- suppose (a,b,c,d,e) is deleted from Movies
- what to delete from MovieProd?
- DELETE FROM MovieProd WHERE title='a'

what if we DELETE (a,b) FROM Producer?

### rewriting queries to use materialized views

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assume a materialized view v defined by

SELECT L<sub>v</sub> FROM R<sub>v</sub> WHERE C<sub>v</sub>

- $L_v$  is a list of attributes
- $\triangleright$   $R_v$  is a list of relations
- $C_v$  is a condition

#### rewriting queries to use materialized views

suppose a query q defined by

SELECT  $L_q$  FROM  $R_q$  WHERE  $C_q$ 

- L<sub>q</sub> is a list of attributes
- $R_q$  is a list of relations
- $C_q$  is a condition

can we rewrite q to use what v materializes?

# conditions for rewriting

if all of the following apply

- 1.  $R_v \subseteq R_q$
- 2.  $C_q \Rightarrow C_v$
- 3. if  $C_q \equiv C_v \wedge C$  then the attributes of relations on  $R_v$  that C mentions are attributes on  $L_v$

4. the attributes on  $L_q$  that come from  $R_v$  are also on  $L_v$ 

# rewriting

if the conditions expressed above are met, then

1. replace  $R_q$  by V and the relations in  $R_q \setminus R_v$ 

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2. replace  $C_q$  by C

#### example

assume the following view is materialized

CREATE MATERIALIZED VIEW MovieProd(movieTitle,prodNat) AS SELECT title,nationality FROM Movies, Producer WHERE producer = name

then the query

SELECT starName,ProdNat FROM Movies, Producer, StarsIn WHERE title=movieTitle AND producer=name

can be rewritten SELECT starName,ProdNat FROM MovieProd, StarsIn WHERE title=movieTitle

#### example

assume the following view is materialized  $MovieProd(t,n) \leftarrow Movies(t,y,g,st,p), Producer(p,n)$  then the query

 $ans(t,n) \leftarrow Movies(t,y,g,st,p), Producer(p,n), StarsIn(t,y,s)$ 

can be rewritten

 $ans(s,n) \leftarrow MovieProd(t,n),StarsIn(t,y,s)$